New in CoCoA-5.2.4 and CoCoALib-0.99570 for SC-Square

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Abstract. CoCoALib is a C++ software library offering operations on polynomials, ideals of polynomials, and related objects. The principal developers of CoCoALib are members of the SC^2 project. We give an overview of the latest developments of the library, especially those relating to the project SC^2 .

The CoCoA software suite includes also the programmable, interactive system CoCoA-5. Most of the operations in CoCoALib are also accessible via CoCoA-5. The programmability of CoCoA-5 together with its interactivity help in fast prototyping and testing conjectures.

1 Introduction

We briefly recall what is described in [2] and [3]. The CoCoA project dates back to 1987, with the first public release of its interactive system in 1989. The main purpose of the project has always been to offer a convenient software laboratory for studying *Computational Commutative Algebra*, especially ideals of multivariate polynomials (e.g. Gröbner bases).

The CoCoA software has a modular design: its "mathematical expertise" resides in the C++ software library [4], while the interactive system [7] uses an interpreter which grants easy access to CoCoALib's capabilities (without requiring any knowledge of C++). All code is free and open source (under licence GPL-3).

We give an overview of the latest developments of the library and of the system (for the summer 2018 release). There are some new aspects of particular interest to the project SC^2 . One feature is an implementation in CoCoALib of a new efficient algorithm for computing the *minimal polynomial* ([5, 3]), and its application to many operations on 0-dimensional ideals — a prototype implementation in CoCoA-5 was mentioned in [3]. In particular, in view of applications to Cylindrical Algebraic Decomposition (CAD), we focus on factoring polynomials over algebraic field extensions, and on evaluating good approximations for their roots.

Another feature of interest to SC^2 is the prototype interface for communication between CoCoALib and MathSAT (Section 4.1).

2 Improving usability of CoCoA for SC²

2.1 Timeout mechanism

A flexible "timeout" mechanism has been added to CoCoALib. It has a simple user interface, and can be used with several function calls to CoCoALib. The exact behaviour of the timeout depends on the specific function: e.g. some functions either return the correct and complete answer within the requested time limit, or throw an exception to say that the result could not be computed quickly enough; other functions, which compute approximations to some exact value, return the best approximation which could be obtained within the given time limit.

One purpose of the timeout mechanism is to allow a "speculative" approach to solving: *i.e.* a potentially costly algorithm may be called with a time limit, and in fortunate cases the answer is returned, but in the worst case of a vain attempt only the specified amount of time has been consumed. For example, one may attempt to find all real solutions to a 0-dimensional polynomial system (internally this computes a Gröbner basis); if successful then the result is valuable, but in some unlucky cases the Gröbner basis computation may be unreasonably slow, so we use the timeout mechanism to avoid the "black hole", and must then proceed without knowing what the solutions to the system are — this technique has already been successfully employed inside CoCoALib itself: for example, in testing primality of zero-dimensional ideals (which is called during the factorization of polynomials over algebraic extensions, Section 3.1).

2.2 Towards Real Solving

We have implemented a prototype for GBasisRealSolveCore which removes some non-real components during Gröbner basis computation. The critical internal operation is a quick function for "approximating" the real radical of a polynomial. In this context "approximating" means applying quick heuristics for determining whether a polynomial admits any real solutions; the heuristic may respond true, false or uncertain. Factors which surely have no real roots are removed, whereas those which do have or may have real roots are retained.

The heuristic employs two other ideas mentioned here: real root counting (via Sturm Sequences), and timeout. First tests show the prototype working well. Further studies are needed: there is a considerable body of work on *real radicals*, where however the emphasis was on returning a complete answer (regardless of overall computational cost) rather than an obtaining a quick "approximation".

3 Algebraic extensions

CoCoALib has had for some time the capability to compute with polynomials whose coefficients lie in an algebraic extension. The extension may be specified by any maximal ideal, and does not require that a primitive element be furnished.

One important step for CAD is the computation of isolating intervals for the real roots of a polynomial with coefficients in a (real) algebraic extension. CoCoALib is not yet able to do this last operation, but we are developing the various components necessary to reach this goal.

3.1 Factorization over algebraic extensions

In this section we describe an effective method for factorizing univariate polynomials over finite field extensions. This method, based on the computation of primary decompositions of zero-dimensional ideals, is described in the PhD thesis [10] of the third author, and took inspiration from [11] and [9]; this work included a full implementation in CoCoA-5 language. It has now been streamlined and ported into CoCoALib-0.99570 (together with all functions derived from the computation of the minimal polynomial — see [5]).

Example 1. Let $L = \mathbb{F}_2[a]/\mathcal{M} \cong \mathbb{F}_8$ be the base field, where $\mathcal{M} = \langle a^3 + a + 1 \rangle$ is a maximal ideal, and consider $f(x) = x^5 + x + 1 \in L[x]$, the polynomial to be factorized.

```
/**/ FF_2 ::= ZZ/(2);
/**/ use S ::= FF_2[a];
/**/ M := ideal(a^3+a+1);
/**/ L := S/M;
/**/ use Lx ::= L[x];
/**/ f := x^5+x+1;
/**/ factor(f);
record[
  RemainingFactor := (1),
  factors := [x+(a^2+a+1), x+(a^2+1), x+(a+1), x^2+x+(1)],
  multiplicities := [1, 1, 1, 1]
]
```

This is the algorithm implemented in CoCoALib. Let $S = K[a_1, \ldots, a_m]$ be a polynomial ring over a perfect field K, let \mathcal{M} be a maximal ideal in S, let $L = S/\mathcal{M}$, and let $f(x) \in L[x]$ be a univariate polynomial. The factorization of f(x) is obtained by computing the primary decomposition of the ideal $\mathcal{M} + \langle f(x) \rangle$ in the polynomial ring $K[a_1, \ldots, a_m, x]$; for the proof of correctness see [10].

Algorithm for Factorization over Algebraic Extension

```
Notation: let S = K[a_1, \ldots, a_m], \mathcal{M} a maximal ideal in S, L = S/\mathcal{M}

Input f(x) polynomial in L[x]

1 create the ring P = K[a_1, \ldots, a_m, x]

let \phi \colon S \longrightarrow P, defined by a_i \mapsto a_i

let \psi \colon P \longrightarrow L[x], defined by a_i \mapsto \bar{a}_i and x \mapsto x

2 let \mathcal{M}P = \langle \phi(g) \mid g \in \text{gens}(\mathcal{M}) \rangle \subseteq P

let f_P \in \psi^{-1}(f), a representative of f in P

let J = \mathcal{M}P + \langle f_P \rangle ideal in P
```

```
3 compute the primary decomposition of J \longrightarrow Q_1 \cap \cdots \cap Q_s

4 compute g_i, the monic generator of \psi(Q_i) in L[x] — note: L[x] is PID

5 compute h_i = \operatorname{rad}(g_i) and let m_i = \frac{\deg(g_i)}{\deg(h_i)}

Output \operatorname{LC}(f) \cdot \prod_{j=1}^s h_i^{m_i}, the factorization of f in L[x]
```

Example 2. The internal computation for Example 1, following the algorithm above, actually performs these steps:

Next we see a call to factor in an algebraic extension which is given by a non-principal, maximal ideal, $L = \mathbb{Q}[a,b]/\langle a^2 - 3, b^2 - ab - 4 \rangle$,

```
/**/ use S ::= QQ[a, b];
/**/ M := ideal(a^2-3, b^2-a*b-4);
/**/ IsMaximal(M); // check that M is a maximal ideal
true
/**/ L := S/M;
/**/ use Lx ::= L[x];
/**/ factor(x^6 +(b^2)*x^4 +(-55*b^2+80)*x^2 +(315*b^2-528));
record[
   RemainingFactor := (1),
   factors := [x^2 +(3*b^2), x +(-b), x +(b)],
   multiplicities := [1, 2, 2]
]
```

3.2 Counting real roots

We have implemented a function for computing a primitive *Sturm sequence* of a given (univariate) polynomial with rational coefficients. Sturm sequences enable one to compute easily the number of real roots a given (univariate) polynomial has in a given real interval; in particular, they can be used to tell whether a given (univariate) polynomial has any real roots. CoCoA also offers a function to count the number of real roots:

```
/**/ W := product([x-k | k in 1..20]);//Wilkinson's polynomial
/**/ W2 := W + (1/7402570310)*x^19;
```

```
/**/ NumRealRoots(W2);
18
/**/ W3 := W + (1/7402570311)*x^19;
/**/ NumRealRoots(W3);
20
```

The interactive CoCoA-5 system offers a suite of interpreted functions for computing tight isolating intervals for real roots. This code is not yet available in CoCoALib, but we are planning to port it (though this lengthy task will likely be completed after the end of this first SC² project). While a list of isolating intervals gives more explicit information than a Sturm sequence, it is also generally more costly to compute.

3.3 Bounds on roots of polynomials

CoCoA now also offers a collection of functions for obtaining a good upper bound for the absolute value of every *complex root* of a given (univariate) polynomial. The main function strikes an automatic balance between speed of computation and tightness of the computed bound; an optional second argument lets the caller choose a different balance. A root bound gives less information than a Sturm sequence but can be computed faster, and may sometimes give enough information to decide unsolvability. A root bound can be computed for any nonconstant (univariate) polynomial, but gives no indication whether any of the roots is real.

```
/**/ H := HermitePoly(50, x); // largest real root is 9.18
/**/ RootBound(f); // fair compromise for speed/accuracy
295/32 // 9.22
/**/ RootBound(f,0); // fastest, 0 iterations
237/8 // 29.625 (quite loose)
/**/ RootBound(f,1); // still fast, 1 iteration
115/8 // 14.375 (better than 0 iters)
```

3.4 Interval arithmetic, and range of a polynomial

We have recently added to CoCoALib a prototype suite of functions for "interval arithmetic" on closed real intervals with rational endpoints. The suite includes a more advanced function for evaluating a given (univariate) polynomial with rational coefficients (or interval coefficients) over a given interval. The result is another interval, comprising effectively a lower bound for the minimum and an upper bound for the maximum for the polynomial over the given interval. In general, the result contains strictly the true interval of reachable values: in fact, the true range interval may have irrational end points; e.g. if $f = x^3 - x$ then its range over the interval [-1,1] has both end points being irrational, namely $[-\frac{2\sqrt{3}}{9},\frac{2\sqrt{3}}{9}]$. Consequently, when using intervals with rational end points only an approximation to the correct answer can be produced.

We are still studying the compromise between speed of computation and tightness of the resulting interval. Naturally, it can help answer some questions of solvability where both the variable and the value of the polynomial are limited to finite intervals. One future aim is to use this suite to allow isolation of real roots of polynomials whose coefficients are real algebraic numbers (each represented as a small isolating interval together with the minimal polynomial).

As an example we can compute the range of values of the 10-th Chebyshev polynomial evaluated on the interval I = [-1,1]. From the theory we know that the true answer is the interval [-1,1]; our current prototype implementation produces a fair approximation, being the wider interval from -1.15 to 1.12. The Chebyshev polynomials are an interesting test case because they have many local maximums and minimums, which become global maximums and minimums when we restrict the domain to the interval I. We do not give an explicit example in CoCoA-5 since the user interface is not yet settled.

4 Interface MatSAT↔CoCoA-5

In [3] we described the first steps in the communication between CoCoALib and MathSAT. Some developments followed, in close collaboration with Alberto Griggio.

The construction in CoCoALib of polynomial rings with user defined orderings (such as elimination orderings) has been reorganized so that it makes fewer repeated checks on the admissibility of the term-ordering. This is not a problem in the usual context of Commutative Algebra, where the cost of the computation of a Gröbner basis exceeds by far the time spent in the preliminary construction of the polynomial ring. But the first collection of examples arising from Math-SAT computations, highlighted that this is indeed an issue when we deal with many indeterminates and relatively simple Gröbner basis computations.

We are now designing an automatic conversion to CoCoA-5/CoCoALib of these examples, providing a wide spectrum of cases to test and develop the new functionality given by GBasisRealSolveCore (see Section 2.2).

The communication between CoCoALib and MathSAT has been improved, so the overhead for the data conversions is now negligible.

We wrote a prototype CoCoA-5 package using MSatLinSolve, the first Math-SAT function in CoCoA-5 (see [3]) for the computation of Gröbner fan as an alternative to the call to GFanRelativeInteriorPoint (part of the communication between CoCoALib and Gfanlib [8]). This is an intriguing application, even though it is not (yet) competitive, because the non-trivial code preparing the input for MSatLinSolve is written in the interpreted language of CoCoA-5.

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